Part IV

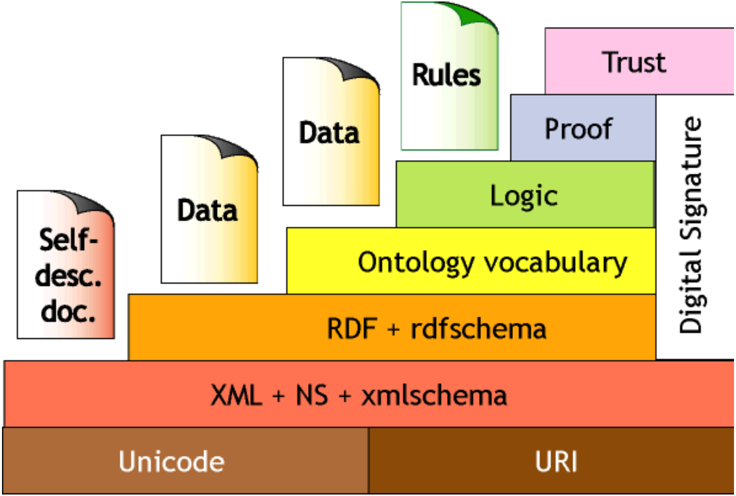
The Semantic Web

Semantic Technologies

- Semantics = meaning
- Semantic Technologies = technologies building on (formalized) meaning
- Declarative Programming as a semantic technology
  - A procedure definition describes its intended meaning
    - e.g. `intersect(L1, L2):- member(X, L1), member(X, L2).`
      Lists `L1` and `L2` intersect if they have a common member `X`.
  - The execution of a program can be viewed as a process of deduction
- The main goal of the Semantic Web (SW) approach:
  - make the information on the web processable by computers
  - machines should be able to understand the web, not only read it
- Achieving the vision of the Semantic Web
  - Add (computer processable) meta-information to the web
  - Formalize background knowledge – build so called ontologies
  - Develop reasoning algorithms and tools
- SW is also used in Knowledge Based Systems (e.g. IBM's Watson)

The vision of the Semantic Web

- The Semantic Web layer cake – Tim Berners-Lee



The Semantic Web

- The goal: making the information on the web processable by computers
- Achieving the vision of the Semantic Web
  - Add meta-information to web pages, e.g.
    (*AIT* `hasLocation` *Budapest*)
    (*AIT* `hasTrack` *Track:Foundational-courses*)
    (*Track:Foundational-courses* `hasCourse` *Semantic-and-declarative...*)
  - Formalise background knowledge – build so called terminologies
    - hierarchies of notions, e.g.
      a *University* is a (subconcept of) *Inst-of-higher-education*,
      the `hasFather` relationship is contained in `hasParent`
    - definitions and axioms, e.g.
      a *Father* is a *Male Person* `having` at least one `child`
  - Develop reasoning algorithms and tools
- Main topics
  - Description Logic, the maths behind the Semantic Web
  - The Web Ontology Language OWL 1 & 2
    from the W3C (WWW Consortium)
  - A glimpse at reasoning algorithms for Description Logic

Contents

First Order Logic (recap)

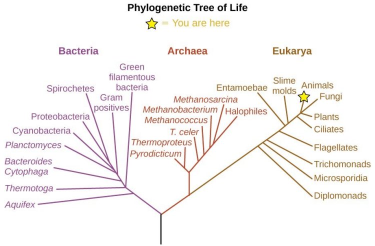
- Syntax:
  - non-logical (“user-defined”) symbols: predicates and functions, including constants (function symbols with 0 arguments)
  - terms (refer to individual elements of the universe, or interpretation), e.g. *fatherOf*(*Susan*)
  - formulas (that hold or do not hold in a given interpretation), e.g. $\varphi = \forall x.(\mathit{Optimist}(\mathit{fatherOf}(x)) \rightarrow \mathit{Optimist}(x))$
- Semantics:
  - determines if a closed formula φ is true in an interpretation $\mathcal{I}$: $\mathcal{I} \models \varphi$ (also read as: $\mathcal{I}$ is a model of φ)
  - an interpretation $\mathcal{I}$ consists of a domain Δ and a mapping from non-logical symbols (e.g. *Optimist*, *fatherOf*, *Susan*) to their meaning
  - semantic consequence: $S \models \alpha$ means: if an interpretation is a model of all formulas in the set S, then it is also a model of α (note that the symbol $\models$ is overloaded)
- Deductive system (also called proof procedure): an algorithm to deduce a consequence α of a set of formulas S: $S \vdash \alpha$
  - example: resolution

Soundness, completeness and decidability (recap)

- A deductive system is **sound** if $S \vdash \alpha \Rightarrow S \models \alpha$ (deduces only truths).
- A deductive system is **complete** if $S \models \alpha \Rightarrow S \vdash \alpha$ (deduces all truths).
- Resolution is a sound and complete deductive system for FOL
- Kurt Gödel was first to show such a system:
  Gödel's completeness theorem: there is a sound and complete deductive system for FOL ($\models \equiv \vdash$)
- FOL is not decidable: no decision procedure for the question "does $S \vdash \alpha$?" (all proofs will be enumerated but no guarantee of obtaining a proof in bounded time – this is also called semi-decidability)
- Developers of the Semantic Web strive for using decidable languages
  - for a decidable language there is a complete and sound deductive system which is guaranteed to terminate within a time limit, which is a function of the (size of) the input formulas
- Semantic Web languages are based on Description Logics, which are decidable sublanguages of FOL

Ontologies

- **Ontology**: computer processable description of knowledge
- Early ontologies include classification system (biology, medicine, books)



- Entities in the Web Ontology Language (OWL):
  - objects – correspond to real life objects (e.g. people, such as Susan, her parents, etc.)
  - classes – describe sets of objects (e.g. optimists)
  - properties (attributes, slots) – describe binary relationships (e.g. has parent)

Knowledge Representation

(Color coding: object–individual, class–concept, property–role)

- **Natural Language**:
  1. someone having a non-optimist friend is bound to be an optimist.
  2. Susan has herself as a friend.
- **First order Logic**:
  1. $\forall x.(\exists y.(hasFriend(x,y) \wedge \neg opt(y)) \rightarrow opt(x))$
  2. $hasFriend(Susan, Susan)$
- **Description Logics**:
  1. $(\exists \mathrm{hasFriend}.\neg \mathrm{Opt}) \sqsubseteq \mathrm{Opt}$ (GCI – gen. concept inclusion axiom)
  2. $\mathrm{hasFriend}(\mathrm{Susan}, \mathrm{Susan})$ (role assertion)
- **Web Ontology Language** (Manchester syntax)[4]:
  1. `(hasFriend some (not Opt)) SubClassOf Opt`
     Those `having some friends` who are `not Opt` must be `Opt`
     (GCI – general class inclusion axiom)
  2. `hasFriend(Susan,Susan)` (object property assertion)

[4] `protegeproject.github.io/protege/class-expression-syntax`

A sample ontology

1. There are “things” that are animals.
2. There are animals that are male.
3. There are animals that are female.
4. There are animals that are humans.
5. There are humans who are optimists.
6. There are animals that are happy.
7. Optimists are always happy.
8. There is a relation “has parent”.
9. There is a relation “has father”, which implies “has parent”.
10. There is a relation “has mother”, which implies “has parent”.
11. The right hand side of “has father” has to be male.
12. The right hand side of “has mother” has to be female.
13. There is a relation “has friend”.
14. Someone having an optimistic parent is optimistic.
15. Someone having a non-optimistic friend is optimistic.
16. There are individuals: Susan, her mother Mother and her father Father.
17. Mother has Father as her friend.

Try drawing conclusions from these statements.

The sample ontology – saved in Description Logic notation

(Select the “save as” format as “Latex syntax” to obtain DL notation.)

2. There are animals that are male. $\text{Male} \sqsubseteq \text{Animal}$
3. There are animals that are female. $\text{Female} \sqsubseteq \text{Animal}$
4. There are animals that are humans. $\text{Human} \sqsubseteq \text{Animal}$
5. There are humans who are optimists $\text{Opt} \sqsubseteq \text{Human}$
6. There are animals that are happy. $\text{Happy} \sqsubseteq \text{Animal}$
7. Optimists are always happy. $\text{Opt} \sqsubseteq \text{Happy}$
9. Relation “has father” implies “has parent”. $\text{hasFather} \sqsubseteq \text{hasParent}$
10. Relation “has mother” implies “has parent”. $\text{hasMother} \sqsubseteq \text{hasParent}$
11. The RHS (range) of “has father” has to be male. $\top \sqsubseteq \forall\ \text{hasFather Male}$
12. The RHS (range) of “has mother” has to be female. $\top \sqsubseteq \forall\ \text{hasMother Female}$
14. Someone with an opt. parent is optimistic. $\text{C1} \equiv \exists\ \text{hasParent Opt},\ \text{C1} \sqsubseteq \text{Opt}$
15. Someone with a non-opt. friend is opt. $\text{C2} \equiv \exists\ \text{hasFriend}\ \neg\text{Opt},\ \text{C2} \sqsubseteq \text{Opt}$
16. There are individuals: Susan, her mother Mother and her father Father.
    hasFather(Susan, Father), hasMother(Susan, Mother)
17. Mother has Father as her friend. hasFriend(Mother, Father)

Description Logic (DL) – an overview

DL, a subset of FOL, is the mathematical background of OWL

- Signature – relation and function symbols allowed in DL
  - role name (R) – binary predicate symbol (cf. OWL property)
  - concept name (A) – unary predicate symbol (cf. OWL class)
  - individual name ($a, \ldots$) – constant symbol (cf. OWL object)
  - No non-constant function symbols, no preds of arity > 2, no vars
- Concept names and concept expressions represent sets, e.g. $\exists$hasParent.Optimist – the set of those who have an optimist parent
- Terminological axioms (TBox) – stating background knowledge
  - A simple axiom using the DL language $\mathcal{ALE}$: $\exists$hasParent.Optimist $\sqsubseteq$ Optimist – the set of those who have an optimist parent is a subset of the set of optimists
  - Translation to FOL: $\forall x.(\exists y.(hasP(x, y) \land Opt(y)) \rightarrow Opt(x))$
- Assertions (ABox) – stating facts about individual names
  - Example: Optimist(JACOB), hasParent(JOSEPH, JACOB)
- A consequence of these TBox and ABox axioms is: Optimist(JOSEPH)
- The Description Logic used in OWL1 and OWL2 are decidable: there are bounded time algorithms for checking the $\models$ relationship

Some further examples of terminological axioms

(1) A Mother is a Person, who is a Female and who has(a)Child.

$$\mathsf{Mother} \equiv \mathsf{Person} \sqcap \mathsf{Female} \sqcap \exists \mathsf{hasChild}.\top$$

(2) A Tiger is a Mammal.

$$\mathsf{Tiger} \sqsubseteq \mathsf{Mammal}$$

(3) Children of an Optimist Person are Optimists, too.

$$\mathsf{Optimist} \sqcap \mathsf{Person} \sqsubseteq \forall \mathsf{hasChild}.\mathsf{Optimist}$$

(4) Childless people are Happy.

$$\forall \mathsf{hasChild}.\bot \sqcap \mathsf{Person} \sqsubseteq \mathsf{Happy}$$

(5) Those in the relation hasChild are also in the relation hasDescendent.

$$\mathsf{hasChild} \sqsubseteq \mathsf{hasDescendent}$$

(6) The relation hasParent is the inverse of the relation hasChild.

$$\mathsf{hasParent} \equiv \mathsf{hasChild}^{-}$$

(7) The hasDescendent relationship is transitive.

$$\mathbf{Trans}(\mathsf{hasDescendent})$$

Description Logics – why the plural?

- These logic variants were progressively developed in the last two decades
- As new constructs were proved to be "safe", i.e. keeping the logic decidable, these were added
- We will start with the very simple language $\mathcal{AL}$, extend it to $\mathcal{ALE}$, $\mathcal{ALU}$ and $\mathcal{ALC}$
- As a side branch we then define $\mathcal{ALCN}$
- We then go back to $\mathcal{ALC}$ and extend it to languages $\mathcal{S}$, $\mathcal{SH}$, $\mathcal{SHI}$ and $\mathcal{SHIQ}$ (which encompasses $\mathcal{ALCN}$)
- We briefly tackle further extensions $\mathcal{O}$, (**D**) and $\mathcal{R}$
- OWL 1, published in 2004, corresponds to $\mathcal{SHOIN}$(**D**)
- OWL 2, published in 2012, corresponds to $\mathcal{SROIQ}$(**D**)

Contents

Overview of the $\mathcal{ALCN}$ language

- In $\mathcal{ALCN}$ a statement (axiom) can be
  - a subsumption (inclusion), e.g. Tiger $\sqsubseteq$ Mammal, or
  - an equivalence, e.g. Woman $\equiv$ Female $\sqcap$ Person, Mother $\equiv$ Woman $\sqcap$ $\exists$hasChild.$\top$
- In general, an $\mathcal{ALCN}$ axiom can take these two forms:
  - subsumption: $C \sqsubseteq D$
  - equivalence: $C \equiv D$, where C and D are concept expressions
- A concept expression C denotes a set of objects (a subset of the Δ universe of the interpretation), and can be:
  - an atomic concept (or concept name), e.g. Tiger, Female, Person
  - a composite concept, e.g. Female $\sqcap$ Person, $\exists$hasChild.$\top$
  - composite concepts are built from atomic concepts and *atomic roles* (also called role names) using some constructors (e.g. $\sqcap$, $\sqcup$, $\exists$, etc.)
- We first introduce language $\mathcal{AL}$, that allows a minimal set of constructors (all examples on this page are valid $\mathcal{AL}$ concept expressions)
- Next, we discuss new constructors named $\mathcal{U}$, $\mathcal{E}$, $\mathcal{C}$, $\mathcal{N}$

The syntax of the $\mathcal{AL}$ language

Language $\mathcal{AL}$ (Attributive Language) allows the following concept expressions, also called concepts, for short:

A is an atomic concept, C, D are arbitrary (possibly composite) concepts
R is an atomic role

DL conc.	OWL class	Name	Informal definition
A	A (class name)	atomic concept	those in A
$\top$	`owl:Thing`	top	the set of all objects
$\bot$	`owl:Nothing`	bottom	the empty set
$\neg A$	`not` A	atomic negation	those not in A
$C \sqcap D$	C `and` D	intersection	those in both C and D
$\forall R.C$	R `only` C	value restriction	those whose all Rs belong to C
$\exists R.\top$	R `some owl:Thing`	limited exist. restr.	those who have at least one R

Examples of $\mathcal{AL}$ concept expressions:

Person $\sqcap$ $\neg$Female	`Person and not Female`
Person $\sqcap$ $\forall$hasChild.Female	`Person and (hasChild only Female)`
Person $\sqcap$ $\exists$hasChild.$\top$	`Person and (hasChild some owl:Thing)`

The semantics of the $\mathcal{AL}$ language

- An interpretation $\mathcal{I}$ is a mapping:
  - $\Delta^\mathcal{I} = \Delta$ is the universe, the **nonempty** set of all individuals/objects
  - for each concept/class name A, $A^\mathcal{I}$ is a (possibly empty) subset of Δ
  - for each role/property name R, $R^\mathcal{I} \subseteq \Delta \times \Delta$ is a bin. relation on Δ
- The semantics of $\mathcal{AL}$ extends $\mathcal{I}$ to composite concept expressions, i.e. describes how to "calculate" the meaning of arbitrary concept exprs

$$
\begin{aligned}
\top^\mathcal{I} &= \Delta \\
\bot^\mathcal{I} &= \emptyset \\
(\neg A)^\mathcal{I} &= \Delta \setminus A^\mathcal{I} \\
(C \sqcap D)^\mathcal{I} &= C^\mathcal{I} \cap D^\mathcal{I} \\
(\forall R.C)^\mathcal{I} &= \{a \in \Delta | \forall b.(\langle a, b \rangle \in R^\mathcal{I} \rightarrow b \in C^\mathcal{I})\} \\
(\exists R.\top)^\mathcal{I} &= \{a \in \Delta | \exists b.\langle a, b \rangle \in R^\mathcal{I}\}
\end{aligned}
$$

- Finally the semantics of axioms maps them to truth values:

$$
\begin{aligned}
\mathcal{I} \models C \sqsubseteq D \quad &\text{iff} \quad C^\mathcal{I} \subseteq D^\mathcal{I} \\
\mathcal{I} \models C \equiv D \quad &\text{iff} \quad C^\mathcal{I} = D^\mathcal{I}
\end{aligned}
$$

The $\mathcal{ALCN}$ language family: extensions $\mathcal{U}$, $\mathcal{E}$, $\mathcal{C}$, $\mathcal{N}$

Further concept constructors, OWL equivalents shown in [`square brackets`]:

- Union: $C \sqcup D$, [C `or` D] – those in either C or D
$$(C \sqcup D)^{\mathcal{I}} = C^{\mathcal{I}} \cup D^{\mathcal{I}} \qquad (\mathcal{U})$$
- Full existential restriction: $\exists R.C$, [R `some` C]
– those who have at least one R belonging to C
$$(\exists R.C)^{\mathcal{I}} = \{a \in \Delta^{\mathcal{I}} | \exists b.\langle a, b \rangle \in R^{\mathcal{I}} \wedge b \in C^{\mathcal{I}}\} \qquad (\mathcal{E})$$
- (Full) negation: $\neg C$, [`not` C] – those who do not belong to C
$$(\neg C)^{\mathcal{I}} = \Delta^{\mathcal{I}} \setminus C^{\mathcal{I}} \qquad (\mathcal{C})$$
- Number restrictions (unqualified): $(\geqslant nR)$, [R `min` n `owl:Thing`] and
$(\leqslant nR)$, [R `max` n `owl:Thing`]
– those who have at least n R-s, or have at most n R-s
$$\begin{aligned}(\geqslant n\,R)^{\mathcal{I}} &= \left\{ a \in \Delta^{\mathcal{I}} \mid \quad \left| \left\{ b \mid \langle a, b \rangle \in R^{\mathcal{I}} \right\} \right| \geq n \right\} \\ (\leqslant n\,R)^{\mathcal{I}} &= \left\{ a \in \Delta^{\mathcal{I}} \mid \quad \left| \left\{ b \mid \langle a, b \rangle \in R^{\mathcal{I}} \right\} \right| \leq n \right\}\end{aligned} \qquad (\mathcal{N})$$
Note that qualified number restrictions, such as $(\geqslant n\,R.C)$ (e.g., those having at least 3 blue-eyed children) are not covered by this extension
- E.g.: $\text{Person} \sqcap ((\leqslant 1\,\text{hasCh}) \sqcup (\geqslant 3\,\text{hasCh})) \sqcap \exists\text{hasCh}.\text{Female}$
`Person and (hasCh max 1 or hasCh min 3) and (hasCh some Female)`

Rewriting $\mathcal{ALCN}$ to first order logic

- Concept expressions map to predicates with one argument, e.g.
  Tiger $\Longrightarrow$ $Tiger(x)$ Mammal $\Longrightarrow$ $Mammal(x)$
  Person $\Longrightarrow$ $Person(x)$ Female $\Longrightarrow$ $Female(x)$
- Simple connectives $\sqcap$, $\sqcup$, $\neg$ map to boolean operations $\wedge$, $\vee$, $\neg$, e.g.
  Person $\sqcap$ Female $\Longrightarrow$ $Person(x) \wedge Female(x)$
  Person $\sqcup$ $\neg$Mammal $\Longrightarrow$ $Person(x) \vee \neg Mammal(x)$
- An axiom $C \sqsubseteq D$ can be rewritten as $\forall x.(C(x) \rightarrow D(x))$, e.g.
  Tiger $\sqsubseteq$ Mammal $\Longrightarrow$ $\forall x.(Tiger(x) \rightarrow Mammal(x))$
- An axiom $C \equiv D$ can be rewritten as $\forall x.(C(x) \leftrightarrow D(x))$, e.g.
  Woman $\equiv$ Person $\sqcap$ Female $\Longrightarrow$ $\forall x.(Woman(x) \leftrightarrow Person(x) \wedge Female(x))$
- Concept constructors involving a role name can be rewritten to a quantified formula.

Rewriting $\mathcal{ALCN}$ to first order logic, example

- Consider $C = \text{Person} \sqcap ((\leqslant 1\,\text{hasCh}) \sqcup (\geqslant 3\,\text{hasCh})) \sqcap \exists\text{hasCh}.\text{Female}$
- Let's outline a predicate $C(x)$ which is true when x belongs to concept C:
  $$C(x) \leftrightarrow Person(x) \wedge (hasAtMost1Child(x) \vee hasAtLeast3Children(x)) \wedge hasFemaleChild(x)$$
- Class practice:
  - Define the FOL predicates $hasAtMost1Child(x)$, $hasAtLeast3Children(x)$, $hasFemaleChild(x)$
  - Additionally, define the following FOL predicates:
    - $hasOnlyFemaleChildren(x)$, corresponding to the concept $\forall\text{hasCh}.\text{Female}$
    - $hasAtMost2Children(x)$, corresponding to the concept $(\leqslant 2\,\text{hasCh})$

Rewriting $\mathcal{ALCN}$ to first order logic, solutions

- $\exists$hasCh.Female
  $hasFemaleChild(x) \leftrightarrow \exists y.(hasCh(x,y) \wedge Female(y))$
- $\forall$hasCh.Female
  $hasOnlyFemaleChildren(x) \leftrightarrow \forall y.(hasCh(x,y) \rightarrow Female(y))$
- ($\leqslant$ 1 hasCh)
  $hasAtMost1Child(x) \leftrightarrow \forall y,z.(hasCh(x,y) \wedge hasCh(x,z) \rightarrow y = z)$
- ($\geqslant$ 3 hasCh)
  $hasAtLeast3Children(x) \leftrightarrow$
  $\exists y,z,w.(hasCh(x,y) \wedge hasCh(x,z) \wedge hasCh(x,w) \wedge y \neq z \wedge y \neq w \wedge z \neq w)$
- ($\leqslant$ 2 hasCh)
  $hasAtMost2Children(x) \leftrightarrow$
  $\forall y,z,w.(hasCh(x,y) \wedge hasCh(x,z) \wedge hasCh(x,w) \rightarrow (y = z \vee y = w \vee z = w))$

General rewrite rules $\mathcal{ALCN} \to$ FOL

Each concept expression can be mapped to a FOL formula:

- Each concept expression C is mapped to a formula $\Phi_C(x)$ (expressing that x belongs to C).
- Atomic concepts (A) and roles (R) are mapped to unary and binary predicates $A(x), R(x, y)$.
- $\sqcap$, $\sqcup$, and $\neg$ are transformed to their counterpart in FOL $(\wedge, \vee, \neg)$, e.g. $\Phi_{C \sqcap D}(x) = \Phi_C(x) \wedge \Phi_D(x)$
- Mapping further concept constructors:

$$\begin{aligned}
\Phi_{\exists R.C}(x) &= \exists y.\,(R(x,y) \wedge \Phi_C(y)) \\
\Phi_{\forall R.C}(x) &= \forall y.\,(R(x,y) \rightarrow \Phi_C(y)) \\
\Phi_{\geqslant n\,R}(x) &= \exists y_1, \ldots, y_n.\left(R(x,y_1) \wedge \cdots \wedge R(x,y_n) \wedge \bigwedge_{i<j} y_i \neq y_j\right) \\
\Phi_{\leqslant n\,R}(x) &= \forall y_1, \ldots, y_{n+1}.\left(R(x,y_1) \wedge \cdots \wedge R(x,y_{n+1}) \rightarrow \bigvee_{i<j} y_i = y_j\right)
\end{aligned}$$

Equivalent languages in the $\mathcal{ALCN}$ family

- Language $\mathcal{AL}$ can be extended by arbitrarily choosing whether to add each of $\mathcal{UECN}$, resulting in $\mathcal{AL}[\mathcal{U}][\mathcal{E}][\mathcal{C}][\mathcal{N}]$.
  Do these $2^4 = 16$ languages have different expressive power?
  Two concept expressions are said to be equivalent, if they have the same meaning, in all interpretations.
  Languages $\mathcal{L}_1$ and $\mathcal{L}_2$ have the same expressive power ($\mathcal{L}_1 \stackrel{e}{=} \mathcal{L}_2$), if any expression of $\mathcal{L}_1$ can be mapped into an equivalent expression of $\mathcal{L}_2$, and vice versa.
- As a preparation for discussing the above let us show that these axioms hold in all models, for arbitrary concepts C and D and role R:

$$
\begin{aligned}
C \sqcup D &\equiv \neg(\neg C \sqcap \neg D) \\
\exists R.C &\equiv \neg \forall R.\neg C
\end{aligned}
$$

$$
\begin{aligned}
\neg\neg C &\equiv C \\
\neg\top &\equiv \bot \\
\neg\bot &\equiv \top \\
\neg(C \sqcap D) &\equiv \neg C \sqcup \neg D \\
\neg\exists R.\top &\equiv \forall R.\bot \\
\neg\forall R.C &\equiv \exists R.\neg C
\end{aligned}
$$

Equivalent languages in the $\mathcal{ALCN}$ family

Let us show that $\mathcal{ALUE}$ and $\mathcal{ALC}$ are equivalent:

- As $C \sqcup D \equiv \neg(\neg C \sqcap \neg D)$ and $\exists R.C \equiv \neg\forall R.\neg C$, union and full existential restriction can be eliminated by using (full) negation. That is, to each $\mathcal{ALUE}$ concept expression there exists an equivalent $\mathcal{ALC}$ expression.
- The other way, each $\mathcal{ALC}$ concept can be transformed to an equivalent $\mathcal{ALUE}$ expression, by moving negation inwards, until before atomic concepts, and removing double negation; using the axioms from the right hand column on the previous slide
- Thus $\mathcal{ALUE}$ and $\mathcal{ALC}$ have the same expressive power, and so $\mathcal{ALC}(\mathcal{N}) \stackrel{e}{=} \mathcal{ALCU}(\mathcal{N}) \stackrel{e}{=} \mathcal{ALCE}(\mathcal{N}) \stackrel{e}{=} \mathcal{ALCUE}(\mathcal{N}) \stackrel{e}{=} \mathcal{ALUE}(\mathcal{N})$.

Further remarks:

- As $\mathcal{U}$ and $\mathcal{E}$ is subsumed by $\mathcal{C}$, we will use $\mathcal{ALC}$ to denote the language allowing $\mathcal{U}$, $\mathcal{E}$ and $\mathcal{C}$
- It can be shown that any two of $\mathcal{AL}, \mathcal{ALU}, \mathcal{ALE}, \mathcal{ALC}, \mathcal{ALN}, \mathcal{ALUN}, \mathcal{ALEN}, \mathcal{ALCN}$ have different expressive power

Contents

A special case of ontology: definitional TBox

- $\mathcal{T}_{fam}$: a sample definitional TBox for family relationships

$$
\begin{aligned}
\mathsf{Woman} &\equiv \mathsf{Person} \sqcap \mathsf{Female} \\
\mathsf{Man} &\equiv \mathsf{Person} \sqcap \neg\mathsf{Woman} \\
\mathsf{Mother} &\equiv \mathsf{Woman} \sqcap \exists\mathsf{hasChild}.\mathsf{Person} \\
\mathsf{Father} &\equiv \mathsf{Man} \sqcap \exists\mathsf{hasChild}.\mathsf{Person} \\
\mathsf{Parent} &\equiv \mathsf{Father} \sqcup \mathsf{Mother} \\
\mathsf{Grandmother} &\equiv \mathsf{Woman} \sqcap \exists\mathsf{hasChild}.\mathsf{Parent}
\end{aligned}
$$

- A definitional TBox consists of equivalence axioms only, the left hand sides being distinct concept names (atomic concepts)
- The concepts on the left hand sides are called name symbols
- The remaining atomic concepts are called base symbols, e.g. in our example the two base symbols are Person and Female.
- In a definitional TBox the meanings of name symbols can be obtained by evaluating the right hand side of their definition

Interpretations and semantic consequence

Recall the definition of assigning a truth value to TBox axioms in an interpretation $\mathcal{I}$:

$$\mathcal{I} \models C \sqsubseteq D \quad \text{iff} \quad C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$$
$$\mathcal{I} \models C \equiv D \quad \text{iff} \quad C^{\mathcal{I}} = D^{\mathcal{I}}$$

Based on this we introduce the notion of “semantic consequence” exactly in the same way as for FOL

- We can naturally extend the above $\mathcal{I} \models \alpha$ notation
  – where α is either $C \sqsubseteq D$ or $C \equiv D$ –
  to a TBox (i.e. a set of α axioms) $\mathcal{T}$
  - $\mathcal{I} \models \mathcal{T}$ ($\mathcal{I}$ satisfies $\mathcal{T}$, $\mathcal{I}$ is a model of $\mathcal{T}$) iff
    for each $\alpha \in \mathcal{T}$, $\mathcal{I} \models \alpha$, i.e. $\mathcal{I}$ is a model of α
- We now overload even further the “ $\models$ ” symbol:
  $\mathcal{T} \models \alpha$ (read axiom α is a semantic consequence of the TBox $\mathcal{T}$) iff
  - all models of $\mathcal{T}$ are also models of α, i.e.
  - for all interpretations $\mathcal{I}$, if $\mathcal{I} \models \mathcal{T}$ holds, then $\mathcal{I} \models \alpha$ also holds

TBox reasoning tasks

Reasoning tasks on TBoxes only (i.e. no ABoxes involved)

- A base assumption: the TBox is **consistent** (does not contain a contradiction), i.e. it has a model
- **Subsumption**: concept C is subsumed by concept D wrt. a TBox $\mathcal{T}$, iff $\mathcal{T} \models (C \sqsubseteq D)$, i.e. $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$ holds in all $\mathcal{I}$ models of $\mathcal{T}$ ($C \sqsubseteq_{\mathcal{T}} D$)
  e.g. $\mathcal{T}_{fam} \models (\text{Grandmother} \sqsubseteq \text{Parent})$ (recall that $\mathcal{T}_{fam}$ is the family TBox)
- **Equivalence**: concepts C and D are equivalent wrt. a TBox $\mathcal{T}$, iff $\mathcal{T} \models (C \equiv D)$, i.e. $C^{\mathcal{I}} = D^{\mathcal{I}}$ holds in all $\mathcal{I}$ models of $\mathcal{T}$ ($C \equiv_{\mathcal{T}} D$).
  e.g. $\mathcal{T}_{fam} \models (\text{Parent} \equiv \text{Person} \sqcap \exists \text{hasChild}.\text{Person})$
- **Disjointness**: concepts C and D are disjoint wrt. a TBox $\mathcal{T}$, iff $\mathcal{T} \models (C \sqcap D \equiv \bot)$, i.e. $C^{\mathcal{I}} \cap D^{\mathcal{I}} = \emptyset$ holds in all $\mathcal{I}$ models of $\mathcal{T}$.
  e.g. $\mathcal{T}_{fam} \models (\text{Woman} \sqcap \text{Man} \equiv \bot$
- Note that all these tasks involve two concepts, C and D

Reducing reasoning tasks to testing satisfiability

- We now introduce a simpler, but somewhat artificial reasoning task: checking the satisfiability of a concept
- **Satisfiability**: a concept C is satisfiable wrt. TBox $\mathcal{T}$, iff there is a model $\mathcal{I}$ of $\mathcal{T}$ such that $C^{\mathcal{I}}$ is non-empty (hence C is non-satisfiable wrt. $\mathcal{T}$ iff in all $\mathcal{I}$ models of $\mathcal{T}$ $C^{\mathcal{I}}$ is empty)
- We will reduce each of the earlier tasks to checking non-satisfiability
- E.g. to prove: Woman $\sqsubseteq$ Person, let's construct a concept C that contains all counter-examples to this statement: $C =$ Woman $\sqcap \neg$Person
- If we can prove that C has to be empty, i.e. there are no counter-examples, then we have proven the subsumption
- Assume we have a method for checking satisfiability. Other tasks can be reduced to this method (usable in $\mathcal{ALC}$ and above):
  - C is subsumed by $D \iff C \sqcap \neg D$ is not satisfiable
  - C and D are equivalent $\iff (C \sqcap \neg D) \sqcup (D \sqcap \neg C)$ is not satisfiable
  - C and D are disjoint $\iff C \sqcap D$ is not satisfiable
- In simpler languages, not supporting full negation, such as $\mathcal{ALN}$, all reasoning tasks can be reduced to subsumption

Contents

The $\mathcal{SHIQ}$ Description Logic language – an overview

- Expanding the abbreviation $\mathcal{SHIQ}$
  - $\mathcal{S} \equiv \mathcal{ALC}_{\mathcal{R}^+}$ (language $\mathcal{ALC}$ extended with transitive roles), i.e. one can state that certain roles (e.g. hasAncestor) are transitive.
  - $\mathcal{H} \equiv$ role hierarchies. Adds statements of the form $R \sqsubseteq S$, e.g. if a pair of objects belongs to the hasFriend relationship, then it must belong to the knows relationship too: hasFriend $\sqsubseteq$ knows (could be stated in English as: *everyone knows their friends*)
  - $\mathcal{I} \equiv$ inverse roles: allows using role expressions R^- to denote the inverse of role R, e.g. hasParent $\equiv$ hasChild$^-$
  - $\mathcal{Q} \equiv$ qualified number restrictions (a generalisation of $\mathcal{N}$): allows the use of concept expressions $(\leqslant nR.C)$ and $(\geqslant nR.C)$ e.g. those who have at least 3 tall children : $(\geqslant 3\,\text{hasChild}.\text{Tall})$

$\mathcal{SHIQ}$ language extensions – the details

- Language $\mathcal{S} \equiv \mathcal{ALC}_{\mathcal{R}^+}$, i.e, $\mathcal{ALC}$ plus transitivity (cf. the index $_{R^+}$)
  - Concept axioms and concept expressions – same as in $\mathcal{ALC}$
  - An additional axiom type: **Trans**(R) declares role R to be transitive
- Extension $\mathcal{H}$ – introducing role hierarchies
  - Adds role axioms of the form $R \sqsubseteq S$ and $R \equiv S$ ($R \equiv S$ can be eliminated, replacing it by $R \sqsubseteq S$ and $S \sqsubseteq R$)
  - In $\mathcal{SH}$ it is possible describe a weak form of transitive closure:

$$\textbf{Trans}(\text{hasDescendent})$$
$$\text{hasChild} \sqsubseteq \text{hasDescendent}$$

  - This means that hasDescendent is a transitive role which includes hasChild
  - What we cannot express in $\mathcal{SH}$ is that hasDescendent is the smallest such role.
    (This property cannot be described in FOL either.)

$\mathcal{SHIQ}$ language extensions – the details (2)

Extension $\mathcal{I}$ – adding inverse roles

- Our first role constructor is $^-$: R^- is the inverse of role R
- Example: consider role axiom $\text{hasChild}^- \equiv \text{hasParent}$ and:

$$\begin{aligned} \text{GoodParent} &\equiv \exists\text{hasChild}.\top \sqcap \forall\text{hasChild}.\text{Happy} \\ \text{MerryChild} &\equiv \exists\text{hasParent}.\text{GoodParent} \end{aligned}$$

  A consequence of the above axioms: $\text{MerryChild} \sqsubseteq \text{Happy}$
- Multiple inverses can be eliminated: $(R^-)^- \equiv R, ((R^-)^-)^- \equiv R^- \ldots.$

$\mathcal{SHIQ}$ language extensions – the details (3)

- Extension $\mathcal{Q}$ – qualified number restrictions – generalizing extension $\mathcal{N}$:
  - $(\leqslant nR.C)$ – the set of those who have at most n R-related individuals belonging to C, e.g.
    $(\leqslant 2 \text{hasChild}.\text{Female})$ – those with at most 2 daughters
  - $(\geqslant nR.C)$ – those with at least n R-related individuals belonging to C
- Important: roles appearing in number restrictions have to be simple.
  (This is because otherwise the decidability of the language would be lost.)
- A role is simple if it is not transitive and does not have a transitive sub-role either
  - Given **Trans**(hasDesc), hasDesc is not simple.
  - If we add further role axioms: hasAnc $\equiv$ hasDesc^{-},
    hasAnc $\sqsubseteq$ hasBloodRelation, then hasBloodRelation is not simple
    - hasAnc is transitive because its inverse hasDesc is such
    - hasBloodRelation has the transitive hasAnc as its sub-role

$\mathcal{SHIQ}$ syntax summary

- Notation
  - A – atomic concept, C, C_i – concept expressions
  - R_A – atomic role, R, R_i – role expressions,
    R_S – simple role expression, with no transitive sub-role
- The syntax of concept expressions

$C \rightarrow$		A	*atomic concept*	$(\mathcal{AL})$
	$\mid$	$\top$	*top – universal concept*	$(\mathcal{AL})$
	$\mid$	$\bot$	*bottom – empty concept*	$(\mathcal{AL})$
	$\mid$	$\neg C$	*negation*	$(\mathcal{C})$
	$\mid$	$C_1 \sqcap C_2$	*intersection*	$(\mathcal{AL})$
	$\mid$	$C_1 \sqcup C_2$	*union*	$(\mathcal{U})$
	$\mid$	$\forall R.C$	*value restriction*	$(\mathcal{AL})$
	$\mid$	$\exists R.C$	*existential restriction*	$(\mathcal{E})$
	$\mid$	$(\geqslant n\, R_S.C)$	*qualified number restriction* (R_S: *simple role)*	$(\mathcal{Q})$
	$\mid$	$(\leqslant n\, R_S.C)$	*qualified number restriction*	$(\mathcal{Q})$

$\mathcal{SHIQ}$ syntax summary (2)

- The syntax of role expressions

$R \rightarrow$		R_A	*atomic role*	$(\mathcal{AL})$
	$\mid$	R^-	*inverse role*	$(\mathcal{I})$

- The syntax of terminological axioms

$T \rightarrow$		$C_1 \equiv C_2$	*concept equivalence axiom*	$(\mathcal{AL})$
	$\mid$	$C_1 \sqsubseteq C_2$	*concept subsumption axiom*	$(\mathcal{AL})$
	$\mid$	$R_1 \equiv R_2$	*role equivalence axiom*	$(\mathcal{H})$
	$\mid$	$R_1 \sqsubseteq R_2$	*role subsumption axiom*	$(\mathcal{H})$
	$\mid$	$\mathbf{Trans}(R)$	*transitivity axiom*	$(\mathcal{R}^+)$

$\mathcal{SHIQ}$ semantics

- The semantics of concept expressions

$$
\begin{array}{rcl}
\top^{\mathcal{I}} & = & \Delta^{\mathcal{I}} \\
\bot^{\mathcal{I}} & = & \emptyset \\
(\neg C)^{\mathcal{I}} & = & \Delta^{\mathcal{I}} \setminus C^{\mathcal{I}} \\
(C_1 \sqcap C_2)^{\mathcal{I}} & = & C_1^{\mathcal{I}} \cap C_2^{\mathcal{I}} \\
(C_1 \sqcup C_2)^{\mathcal{I}} & = & C_1^{\mathcal{I}} \cup C_2^{\mathcal{I}} \\
(\forall R.C)^{\mathcal{I}} & = & \{a \in \Delta^{\mathcal{I}} \mid \quad \forall b.\langle a, b \rangle \in R^{\mathcal{I}} \to b \in C^{\mathcal{I}}\} \\
(\exists R.C)^{\mathcal{I}} & = & \{a \in \Delta^{\mathcal{I}} \mid \quad \exists b.\langle a, b \rangle \in R^{\mathcal{I}} \land b \in C^{\mathcal{I}}\} \\
(\geqslant n\, R.C)^{\mathcal{I}} & = & \{\, a \in \Delta^{\mathcal{I}} \mid \quad |\{b \mid \langle a, b \rangle \in R^{\mathcal{I}} \land b \in C^{\mathcal{I}}\}| \geq n\,\} \\
(\leqslant n\, R.C)^{\mathcal{I}} & = & \{\, a \in \Delta^{\mathcal{I}} \mid \quad |\{b \mid \langle a, b \rangle \in R^{\mathcal{I}} \land b \in C^{\mathcal{I}}\}| \leq n\,\}
\end{array}
$$

- The semantics of role expressions

$$
(R^{-})^{\mathcal{I}} \quad = \quad \{\langle b, a \rangle \in \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \mid \langle a, b \rangle \in R^{\mathcal{I}}\}
$$

$\mathcal{SHIQ}$ semantics (2)

- The semantics of terminological axioms

$$
\begin{array}{lcl}
\mathcal{I} \models C_1 \equiv C_2 & \Leftrightarrow & C_1^{\mathcal{I}} = C_2^{\mathcal{I}} \\
\mathcal{I} \models C_1 \sqsubseteq C_2 & \Leftrightarrow & C_1^{\mathcal{I}} \subseteq C_2^{\mathcal{I}} \\
\mathcal{I} \models R_1 \equiv R_2 & \Leftrightarrow & R_1^{\mathcal{I}} = R_2^{\mathcal{I}} \\
\mathcal{I} \models R_1 \sqsubseteq R_2 & \Leftrightarrow & R_1^{\mathcal{I}} \subseteq R_2^{\mathcal{I}} \\
\mathcal{I} \models \mathbf{Trans}(R) & \Leftrightarrow & (\forall a, b, c \in \Delta^{\mathcal{I}}) \\
 & & (\langle a, b \rangle \in R^{\mathcal{I}} \wedge \langle b, c \rangle \in R^{\mathcal{I}} \rightarrow \langle a, c \rangle \in R^{\mathcal{I}})
\end{array}
$$

- Read $\mathcal{I} \models T$ as: “$\mathcal{I}$ satisfies axiom T” or as “$\mathcal{I}$ is a model of T”

Negation normal form (NNF)

- Various normal forms are used in reasoning algorithms
- The tableau algorithms use NNF: only atomic negation allowed
- To obtain NNF, apply the following rules to subterms repeatedly while a subterm matching a left hand side can be found:

$$
\begin{array}{rcl}
\neg\neg C & \rightsquigarrow & C \\
\neg(C \sqcap D) & \rightsquigarrow & \neg C \sqcup \neg D \\
\neg(C \sqcup D) & \rightsquigarrow & \neg C \sqcap \neg D \\
\neg(\exists R.C) & \rightsquigarrow & \forall R.(\neg C) \\
\neg(\forall R.C) & \rightsquigarrow & \exists R.(\neg C) \\
\neg(\leqslant nR.C) & \rightsquigarrow & (\geqslant kR.C) \text{ where } k = n+1 \\
\neg(\geqslant 1R.C) & \rightsquigarrow & \forall R.(\neg C) \\
\neg(\geqslant nR.C) & \rightsquigarrow & (\leqslant kR.C) \text{ if } n > 1, \text{ where } k = n-1
\end{array}
$$

Going beyond $\mathcal{SHIQ}$

- Extension $\mathcal{O}$ introduces nominals, i.e. concepts which can only have a single element. Example: {EUROPE} is a concept whose interpretation must contain a single element
  FullyEuropean $\equiv \forall$hasSite.$\forall$hasLocation.{EUROPE}
- Extension (**D**): concrete domains, e.g. integers, strings etc, whose interpretation is fixed, cf. data properties in OWL
- The Web Ontology Language OWL 1 implements $\mathcal{SHOIN}$(**D**)
- OWL 2 implements $\mathcal{SROIQ}$(**D**)
- The main novelty in $\mathcal{R}$ wrt. $\mathcal{H}$ is the possibility to use role composition ($\circ$):
  hasParent $\circ$ hasBrother $\sqsubseteq$ hasUncle
  i.e. one's parent's brother is one's uncle
- To ensure decidability, the use of role composition is seriously restricted (e.g. it is not allowed to have $\equiv$ instead of $\sqsubseteq$ in the above example)

Contents

The notion of ABox

- The ABox contains assertions about individuals, referred to by individual names a, b, c etc.
  Convention: concrete individual names are written in ALL_CAPITALS
  - concept assertions: $C(a)$, e.g. Father(ALEX), $(\exists$hasJob$.\top)$(BOB)
  - role assertions: $R(a, b)$, e.g. hasChild(ALEX, BOB).
- Individual names correspond to constant symbols of first order logic
- The interpretation function has to be extended:
  - to each individual name a, $\mathcal{I}$ assigns $a^{\mathcal{I}} \in \Delta^{\mathcal{I}}$
- The semantics of ABox assertions is straightforward:
  - $\mathcal{I}$ satisfies a concept assertion $C(a)$ $(\mathcal{I} \models C(a))$, iff $a^{\mathcal{I}} \in C^{\mathcal{I}}$,
  - $\mathcal{I}$ satisfies a role assertion $R(a, b)$ $(\mathcal{I} \models R(a, b))$, iff $\langle a^{\mathcal{I}}, b^{\mathcal{I}} \rangle \in R^{\mathcal{I}}$,
  - $\mathcal{I}$ satisfies an ABox $\mathcal{A}$ $(\mathcal{I} \models \mathcal{A})$ iff $\mathcal{I}$ satisfies all assertions in $\mathcal{A}$, i.e. for all $\alpha \in \mathcal{A}$, $\mathcal{I} \models \alpha$ holds

Reasoning on ABoxes

- ABox $\mathcal{A}$ is consistent wrt. TBox $\mathcal{T}$ iff there is a model $\mathcal{I}$ that satisfies both $\mathcal{A}$ and $\mathcal{T}$ ($\mathcal{I} \models \mathcal{A}$ and $\mathcal{I} \models \mathcal{T}$)
  - Is the ABox $\{\mathsf{Mother}(\mathsf{S}), \mathsf{Father}(\mathsf{S})\}$ consistent wrt. an empty TBox? And wrt. the family TBox?
- Assertion α is a consequence of the ABox $\mathcal{A}$ wrt. TBox $\mathcal{T}$ ($\mathcal{A} \models_{\mathcal{T}} \alpha$) iff $\mathcal{I} \models \alpha$ holds for any interpretation $\mathcal{I}$ for which both $\mathcal{I} \models \mathcal{A}$ and $\mathcal{I} \models \mathcal{T}$ hold
- Example: let $\mathcal{T}$ refer to the family TBox, and $\mathcal{A}$ to the ABox below:

  hasChild(SAM, SUE)   Person(SAM)   Person(SUE)   Person(ANN)
  hasChild(SUE, ANN)   Female(SUE)   Female(ANN)

  Which of the assertions below is a consequence of $\mathcal{A}$ wrt. $\mathcal{T}$?

  1. Mother(SUE)
  2. Mother(SAM)
  3. $\neg$Mother(SAM)
  4. Father(SAM)
  5. (Mother$\sqcup$Father)(SAM)
  6. ($\leqslant$ 1 hasChild)(SAM)

ABoxes and databases

- An ABox may seem similar to a relational database, but
  - Querying a database uses the closed world assumption (CWA): is the query true in the world (interpretation) where the given and only given facts hold?
  - Contrastingly, ABox reasoning uses logical consequence, also called open world assumption (OWA): is it the case that the query holds in all interpretations satisfying the given facts
- At first one may think that with CWA one can always get more deduction possibilities
- However, case-based reasoning in OWA can lead to deductions not possible with CWA (e.g. Susan being optimistic)

Some important ABox reasoning tasks

- Instance check: Decide if assertion α is a consequence of ABox $\mathcal{A}$ wrt. $\mathcal{T}$. Example:
  Check if Mother(SUE) holds wrt. the example ABox $\mathcal{A}$ (on the previous but one slide) and the family TBox.
- Instance retrieval:
  Given a concept expression C find the set of all individual names x such that $\mathcal{A} \models_{\mathcal{T}} C(x)$
  Example: Find all individual names known to belong to the concept Mother

The optimists example as an ABox reasoning task

- Our earlier example of optimists:
  1. If someone has an optimistic parent, then she is optimistic herself.
  2. If someone has a non-optimistic friend, then she is optimistic.
  3. Susan's maternal grandfather has her maternal grandmother as a friend.
- Consider the following TBox $\mathcal{T}$:

$$\exists \mathsf{hP}.\mathsf{Opt} \sqsubseteq \mathsf{Opt} \tag{1}$$

$$\exists \mathsf{hF}.\neg\mathsf{Opt} \sqsubseteq \mathsf{Opt} \tag{2}$$

- Consider the following ABox $\mathcal{A}$, representing (3):
  $\mathsf{hP}(\mathsf{S}, \mathsf{SM})$ $\quad$ $\mathsf{hP}(\mathsf{SM}, \mathsf{SMM})$ $\quad$ $\mathsf{hP}(\mathsf{SM}, \mathsf{SMF})$ $\quad$ $\mathsf{hF}(\mathsf{SMF}, \mathsf{SMM})$
- An instance retrieval task: find the set of all individual names x such that $\mathcal{A} \models_{\mathcal{T}} \mathsf{Opt}(x)$

Another classical example requiring case analysis

- Some facts about the Oedipus family (ABox $\mathcal{A}_{OE}$):

  hasChild(IOCASTE, OEDIPUS)
  hasChild(IOCASTE, POLYNEIKES)
  hasChild(OEDIPUS, POLYNEIKES)
  hasChild(POLYNEIKES, THERSANDROS)

  Patricide(OEDIPUS)

  (¬Patricide)(THERSANDROS)

- Let us call a person "special" if they have a child who is a patricide and who, in turn, has a child who is not a patricide:

$$\text{Special} \equiv \exists\text{hasChild}.(\text{Patricide} \sqcap \exists\text{hasChild}.\neg\text{Patricide})$$

- Let TBox $\mathcal{T}_{OE}$ contain the above axiom only.
- Consider the instance check "Is Iocaste special?":
  $\mathcal{A}_{OE} \models_{\mathcal{T}_{OE}} \text{Special}(\text{IOCASTE})$?
- The answer is "yes", but proving this requires case analysis

Contents

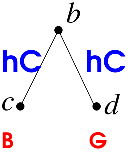
Tableau algorithms

- Various TBox and ABox reasoning tasks have been presented earlier
- In $\mathcal{ALC}$ and above, any TBox task can be reduced to checking satisfiability
- Principles of the $\mathcal{ALCN}$ tableau algorithm
  - It checks if a concept is satisfiable, by trying to construct a model
  - Uses NNF, i.e. "$\neg$" can appear only in front of atomic concepts
  - The model is built through a series of transformations
- The data structure representing the model is called the tableau (state):
  - a directed graph
  - the vertices can be viewed as the domain of the interpretation
  - edges correspond to roles, each edge is labelled by a role
  - vertices are labelled with sets of concepts, to which the vertex is expected to belong
- Example: If a person has a green-eyed and a blonde child, does it follow that she/he has to have a child who is both green-eyed and blonde?
- Formalize the above task as a question in the Description Logic $\mathcal{ALC}$: Does the axiom $(\exists hC.B) \sqcap (\exists hC.G) \sqsubseteq \exists hC.(B \sqcap G)$ hold?[5]

[5] (hC = has child, B = blonde, G = green-eyed)

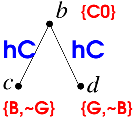
An introductory example, using $\mathcal{ALC}$

- Question: Does the axiom $(\exists \text{hC.B}) \sqcap (\exists \text{hC.G}) \sqsubseteq \exists \text{hC.}(\text{B} \sqcap \text{G})$ hold? (1)
- Reformulate: "Is C satisfiable?", $C = (\exists \text{hC.B}) \sqcap (\exists \text{hC.G}) \sqcap \neg(\exists \text{hC.}(\text{B} \sqcap \text{G}))$ ($A \sqsubseteq B \Leftrightarrow A \sqcap \neg B$ *not satisfiable*)
- The neg. normal form of C is: $C_0 = (\exists \text{hC.B}) \sqcap (\exists \text{hC.G}) \sqcap \forall \text{hC.}(\neg \text{B} \sqcup \neg \text{G}))$
- Goal: build an interpretation $\mathcal{I}$ such that $C_0^{\mathcal{I}} \neq \emptyset$. Thus we try to have a b such that $b \in (\exists \text{hC.B})^{\mathcal{I}}$, $b \in (\exists \text{hC.G})^{\mathcal{I}}$, and $b \in (\forall \text{hC.}(\neg \text{B} \sqcup \neg \text{G}))^{\mathcal{I}}$.
- From $b \in (\exists \text{hC.B})^{\mathcal{I}} \implies \exists c$ such that $\langle b, c \rangle \in \text{hC}^{\mathcal{I}}$ and $c \in \text{B}^{\mathcal{I}}$.
Similarly, $b \in (\exists \text{hC.G})^{\mathcal{I}} \implies \exists d$, such that $\langle b, d \rangle \in \text{hC}^{\mathcal{I}}$ and $d \in \text{G}^{\mathcal{I}}$.
- As b belongs to $\forall \text{hC.}(\neg \text{B} \sqcup \neg \text{G})$, and both c and d are hC relations of b, we obtain constraints: $c \in (\neg \text{B} \sqcup \neg \text{G})^{\mathcal{I}}$ and $d \in (\neg \text{B} \sqcup \neg \text{G})^{\mathcal{I}}$.
- $c \in (\neg \text{B} \sqcup \neg \text{G})^{\mathcal{I}}$ means that either $c \in (\neg \text{B})^{\mathcal{I}}$ or $c \in (\neg \text{G})^{\mathcal{I}}$. Assuming $c \in (\neg \text{B})^{\mathcal{I}}$ contradicts $c \in \text{B}^{\mathcal{I}}$. Thus we have to choose the option $c \in (\neg \text{G})^{\mathcal{I}}$. Similarly, we obtain $d \in (\neg \text{B})^{\mathcal{I}}$.
- We arrive at: $\Delta^{\mathcal{I}} = \{b, c, d\}$;
$\text{hC}^{\mathcal{I}} = \{\langle b, c \rangle, \langle b, d \rangle\}$;
$\text{B}^{\mathcal{I}} = \{c\}$ and $\text{G}^{\mathcal{I}} = \{d\}$.
Here $b \in C_0^{\mathcal{I}}$, thus (1) does not hold.



Extending the example to $\mathcal{ALCN}$

- Question: If a person having at most one child has a green-eyed and a blonde child, does it follow that she/he has to have a child who is both green-eyed and blonde?
- DL question: $(\leqslant 1\text{hC}) \sqcap (\exists\text{hC.B}) \sqcap (\exists\text{hC.G}) \overset{?}{\sqsubseteq} \exists\text{hC.}(\text{B} \sqcap \text{G}))$ (2)
- Reformulation: "Is C satisfiable?", where $C = (\leqslant 1\text{hC}) \sqcap (\exists\text{hC.B}) \sqcap (\exists\text{hC.G}) \sqcap \neg(\exists\text{hC.}(\text{B} \sqcap \text{G}))$
- Negation normal form: $C_0 = (\leqslant 1\text{hC}) \sqcap (\exists\text{hC.B}) \sqcap (\exists\text{hC.G}) \sqcap \forall\text{hC.}(\neg\text{B} \sqcup \neg\text{G}))$
- We first build the same tableau as for (1):



- From $(\leqslant 1\text{hC})(b)$, $\text{hC}(b, c)$, and $\text{hC}(b, d)$ it follows that $c = d$ has to be the case. However merging c and d results in an object being both B and $\neg$B which is a contradiction (clash)
- Thus we have shown that C_0 cannot be satisfied, and thus the answer to question (2) is yes.

Contents

The $\mathcal{ALCN}$ tableau algorithm for empty TBoxes – outline

- “Is C satisfiable?” $\Longrightarrow$ Let’s build a model satisfying C, exhaustively.
- First, bring C to negation normal form C_0.
- The main data structure, the tableau structure $T = (V, E, \mathcal{L}, I)$ where $(V, E, \mathcal{L})$ is a finite directed graph (more about I later)
  - Nodes of the graph (V) can be thought of as domain elements.
  - Edges of the graph (E) represent role relationships between nodes.
  - The labeling function $\mathcal{L}$ assigns labels to nodes and edges:
    - $\forall x \in V$, $\mathcal{L}(x) \subseteq sub(C_0)$, the set of subexpressions of C_0
    - $\forall \langle x, y \rangle \in E$, $\mathcal{L}(\langle x, y \rangle)$ is a role within C (in $\mathcal{SHIQ}$: set of roles)
  - The initial tableau has a single node, the root: $(\{x_0\}, \emptyset, \mathcal{L}, \emptyset)$, where $\mathcal{L}(x_0) = \{C_0\}$. Here C_0 is called the root concept.
- The algorithm uses transformation rules to extend the tableau
- Certain rules are nondeterministic, creating a choice point; backtracking occurs when a trivial clash appears (e.g. both A and $\neg A \in \mathcal{L}(x)$)
- If a clash-free and complete tableau (no rule can fire) is reached $\Longrightarrow$ C is satisfiable.
- When the whole search tree is traversed $\Longrightarrow$ C is not satisfiable.

Outline of the $\mathcal{ALCN}$ tableau algorithm (2)

- The tableau tree is built downwards from the root (edges are always directed downwards)
  - A node b is called an R-successor (or simply successor) of a iff there is an edge from a to b with R as its label, i.e. $\mathcal{L}(\langle a, b \rangle) = R$
- Handling equalities and inequalities
  - To handle $(\leqslant n\,R)$ we need to merge (identify) nodes
  - In handling $(\geqslant n\,R)$ we will have to introduce n R-successors which are pairwise non-identifiable ($x \dot{\neq} y$: x and y are not identifiable)
  - The component I of the tableau data structure $T = (V, E, \mathcal{L}, I)$ is a set of inequalities of the form $x \dot{\neq} y$

Transformation rules of the $\mathcal{ALCN}$ tableau algorithm (1)

$\sqcap$-rule

Condition: $(C_1 \sqcap C_2) \in \mathcal{L}(x)$ and $\{C_1, C_2\} \not\subseteq \mathcal{L}(x)$

New state* T$'$*: $\mathcal{L}'(x) = \mathcal{L}(x) \cup \{C_1, C_2\}$.

$\sqcup$-rule

Condition: $(C_1 \sqcup C_2) \in \mathcal{L}(x)$ and $\{C_1, C_2\} \cap \mathcal{L}(x) = \emptyset$.

New state* T$_1$*: $\mathcal{L}'(x) = \mathcal{L}(x) \cup \{C_1\}$.

New state* T$_2$*: $\mathcal{L}'(x) = \mathcal{L}(x) \cup \{C_2\}$.

$\exists$-rule

Condition: $(\exists R.C) \in \mathcal{L}(x)$, x has no R-successor y s.t. $C \in \mathcal{L}(y)$.

New state* T$'$*: $V' = V \cup \{y\}$ (y is a new node),

$E' = E \cup \{\langle x, y \rangle\}$, $\mathcal{L}'(\langle x, y \rangle) = R$, $\mathcal{L}'(y) = \{C\}$.

$\forall$-rule

Condition: $(\forall R.C) \in \mathcal{L}(x)$, x has an R-successor y s.t. $C \notin \mathcal{L}(y)$.

New state* T$'$*: $\mathcal{L}'(y) = \mathcal{L}(y) \cup \{C\}$.

Transformation rules of the $\mathcal{ALCN}$ tableau algorithm (2)

$\geqslant$-rule

Condition: $(\geqslant n\, R) \in \mathcal{L}(x)$ and x has no n R-successors such that any two are non-identifiable.

New state* $\mathbf{T}'$*:
$V' = V \cup \{y_1, \ldots, y_n\}$ (y_i new nodes),
$E' = E \cup \{\langle\, x, y_1 \,\rangle, \ldots, \langle\, x, y_n \,\rangle\}$,
$\mathcal{L}'(\langle\, x, y_i \,\rangle) = R$, $\mathcal{L}'(y_i) = \emptyset$, for each $i = 1 \leq i \leq n$,
$I' = I \cup \{y_i \dot{\neq} y_j \mid 1 \leq i < j \leq n\}$.

Transformation rules of the $\mathcal{ALCN}$ tableau algorithm (3)

$\leqslant$-rule

Condition: $(\leqslant\ n\,R) \in \mathcal{L}(x)$ and x has R-successors $y_1, \ldots, y_{n+1}$ among which there are at least two identifiable nodes.

For each i and j, $1 \leq i < j \leq n+1$, where y_i and y_j are identifiable:

New state $\mathbf{T}_{ij}$***:*** $V' = V \setminus \{y_j\}$, $\mathcal{L}'(y_i) = \mathcal{L}(y_i) \cup \mathcal{L}(y_j)$,

$E' = E \setminus \{\langle\, x, y_j \,\rangle\} \setminus \{\langle\, y_j, u \,\rangle | \langle\, y_j, u \,\rangle \in E\} \ \cup$

$\{\langle\, y_i, u \,\rangle | \langle\, y_j, u \,\rangle \in E\}$,

$\mathcal{L}'(\langle\, y_i, u \,\rangle) = \mathcal{L}(\langle\, y_j, u \,\rangle)$, $\forall u$ such that $\langle\, y_j, u \,\rangle \in E$,

$I' = I[y_j \to y_i]$ (every occurrence of y_j is replaced by y_i).

The $\mathcal{ALCN}$ tableau algorithm – further details

- There is clash at some node x of a tableau state iff
  - $\{\bot\} \subseteq \mathcal{L}(x)$; or
  - $\{A, \neg A\} \subseteq \mathcal{L}(x)$ for some atomic concept A; or
  - $(\leqslant nR) \in \mathcal{L}(x)$ and x has R-successors $y_1, \ldots, y_{n+1}$ where for any two successors y_i and y_j it holds that $y_i \dot{\neq} y_j \in I$.
- A tableau state is said to be complete, if no transformation rules can be applied at this state (there is no rule the conditions of which are satisfied)

The $\mathcal{ALCN}$ tableau algorithm

In this version the algorithm handles a set of tableau states, one for each yet unexplored subtree of the search space.

1. Intialise the variable `States` = $\{\mathbf{T}_0\}$ (a singleton set containing the initial tableau state)
2. If there is $\mathbf{T} \in$ `States` such that **T** contains a clash, remove **T** from `States` and continue at step 2
3. If there is $\mathbf{T} \in$ `States` such that **T** is complete (and clash-free), exit the algorithm, reporting satisfiability
4. If `States` is empty, exit the algorithm, reporting non-satisfiability
5. Choose an arbitrary element $\mathbf{T} \in$ `States` and apply to **T** an arbitrary transformation rule, whose conditions are satisfied[6] (don't care nondeterminism). Remove **T** from `States`, and add to `States` the `NewStates` resulting from the applied transformation, where `NewStates` = $\{\mathbf{T_1}, \mathbf{T_2}\}$ for the $\sqcup$-**rule**, `NewStates` = $\{\mathbf{T_{ij}} | \cdots\}$ for the $\leqslant$-**rule**, and `NewStates` = $\{\mathbf{T}'\}$ for all other (deterministic) rules. Continue at step 2

[6] Such a tableau state **T** and such a rule exist, because `States` is nonempty, and none of its elements is a complete tableau

The $\mathcal{ALCN}$ tableau algorithm – an example

- Consider checking the satisfiability of concept C_0 (hC = has child, B = blonde):

$$
\begin{aligned}
C_0 &= C_1 \sqcap C_2 \sqcap C_3 \sqcap C_4 \\
C_1 &= (\geqslant 2\,\mathsf{hC}) \\
C_2 &= \exists \mathsf{hC}.\mathsf{B} \\
C_3 &= (\leqslant 2\,\mathsf{hC}) \\
C_4 &= C_5 \sqcup C_6 \\
C_5 &= \forall \mathsf{hC}.\neg \mathsf{B} \\
C_6 &= \mathsf{B}
\end{aligned}
$$

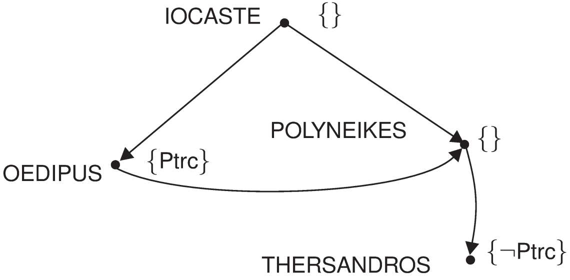
- The tableau algorithm completes with the answer: concept C_0 is satisfiable
- The interpretation constructed by the tableau algorithm: $\Delta^{\mathcal{I}} = \{b, c, d\}$; $\mathsf{hC}^{\mathcal{I}} = \{\langle b, c \rangle, \langle b, d \rangle\}$; $\mathsf{B}^{\mathcal{I}} = \{b, c\}$

Extending the tableau algorithm to ABox reasoning

- To solve an ABox reasoning task (with no TBox), we transform the ABox to a graph, serving as the initial tableau state, e.g. for the IOCASTE family ABox:

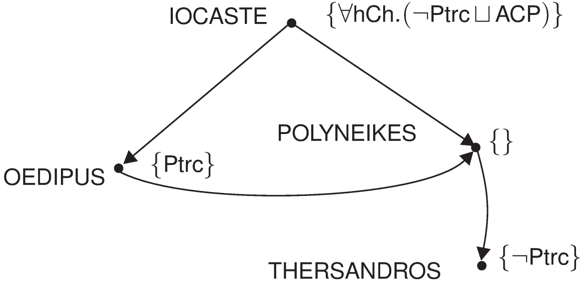
hC(IOCASTE,OEDIPUS)	hC(IOCASTE,POLYNEIKES)
hC(OEDIPUS,POLYNEIKES)	hC(POLYNEIKES,THERSANDROS)
Ptrc(OEDIPUS)	(¬ Ptrc)(THERSANDROS)

- Individual names become nodes of the graph, labelled by a set of concepts, and each role assertion generates an edge, labelled (implicitly) by hC:



Handling ABox axioms in the tableau algorithm (ctd.)

- Given the Iocaste ABox, we want to prove that IOCASTE is special, i.e. she belongs to the concept $\exists hC.(Ptrc \sqcap \exists hC.\neg Ptrc)$
- We do an indirect proof: assume that IOCASTE is not special, i.e. IOCASTE belongs to $(\forall hC.(\neg Ptrc \sqcup \forall hC.Ptrc))$ (1)
- Let's introduce an abbreviation: $ACP \equiv \forall hC.Ptrc$
- To prove that Iocaste is special, we add concept (1) to the IOCASTE node:



- The tableau algorithm, with this initial state, will detect non-satisfiability

Handling TBox axioms in the tableau algorithm

- An arbitrary $\mathcal{ALCN}$ TBox can be transformed to a set of subsumptions of the form $C \sqsubseteq D$ ($C \equiv D$ can be replaced by $\{C \sqsubseteq D, D \sqsubseteq C\}$)
- $C \sqsubseteq D$ can be replaced by $\top \sqsubseteq \neg C \sqcup D$
  cf. $(\alpha \rightarrow \beta)$ is the same as $(\neg\alpha \vee \beta)$
- An arbitrary TBox $\{C_1 \sqsubseteq D_1, C_2 \sqsubseteq D_2, \ldots, C_n \sqsubseteq D_n\}$ can be transformed to a single equivalent axiom: $\top \sqsubseteq C_{\mathcal{T}}$, where

$$C_{\mathcal{T}} = (\neg C_1 \sqcup D_1) \sqcap (\neg C_2 \sqcup D_2) \sqcap \cdots \sqcap (\neg C_n \sqcup D_n).$$

- Concept $C_{\mathcal{T}}$ is called the internalisation of TBox $\mathcal{T}$
- An interpretation $\mathcal{I}$ is a model of a TBox $\mathcal{T}$ ($\mathcal{I} \models \mathcal{T}$) iff each element of the domain belongs to the $C_{\mathcal{T}}$ internalisation concept
  - This observation can be used in the tableaux reasoning algorithm, which tries to build a model
  - To build a model which satisfies the TBox $\mathcal{T}$ we add the concept $C_{\mathcal{T}}$ to the label of each node of the tableau

Handling TBoxes in the tableau algorithm – problems

- Example: Consider the task of checking the satisfiability of concept Blonde wrt. TBox $\{\top \sqsubseteq \exists$hasFriend.Blonde$\}$
  - Concept $\exists$hasFriend.Blonde will appear in each node
  - thus the $\exists$-rule will generate an infinite chain of hasFriend successors
- To prevent the algorithm from looping the notion of blocking is introduced.

Blocking

- Definition: Node y is blocked by node x, if y is a descendant of x and the blocking condition $\mathcal{L}(y) \subseteq \mathcal{L}(x)$ holds (*subset blocking*).
- When y is blocked, we disallow generator rules ($\exists$- and $\geqslant$-rules, creating new successors for y)
- This solves the termination problem, but raises the following issue
  - How can one get an interpretation from the tableau?
  - Solution (approximation, for $\mathcal{ALC}$ only): identify blocked node y with blocking node x, i.e. redirect the edge pointing to y so that it points to x. This creates a model, as
    - all concepts in the label of y are also present in x
    - thus x belongs to all concepts y is expected to belong to
- Is Happy $\sqcap$ Blonde satisfiable wrt. TBox $\{\top \sqsubseteq \exists$hasFriend.Blonde$\}$?

$$\begin{array}{rcl} & x \ \circ & \{\text{Happy}, \text{Blonde}, \exists\text{hasFriend}.\text{Blonde}\} \\ \text{hasFriend} & | & \\ & y \ \circ & \{\text{Blonde}, \exists\text{hasFriend}.\text{Blonde}\} \end{array}$$

  - x blocks y, the tableau is clash-free and complete
  - The model:
    $\Delta^{\mathcal{I}} = \{x\}$; $\text{Happy}^{\mathcal{I}} = \{x\}$; $\text{Blonde}^{\mathcal{I}} = \{x\}$; $\text{hasFriend}^{\mathcal{I}} = \{\langle x, x \rangle\}$